

Post's Correspondence  
Problem

&

Tilings

# PCP

set of dominos  $\left[ \begin{array}{c} \alpha \\ \beta \end{array} \right] \leftarrow 2 \text{ strings}$

decision problem  
given set of dominos

$$S = \left\{ \left[ \begin{array}{c} b \\ ca \end{array} \right], \left[ \begin{array}{c} a \\ ab \end{array} \right], \left[ \begin{array}{c} ca \\ a \end{array} \right], \left[ \begin{array}{c} abc \\ c \end{array} \right] \right\}$$

is there a sequence (with repeats)  
sit. top string = bottom string

$$\left[ \begin{array}{c} a \\ ab \end{array} \right] \left[ \begin{array}{c} b \\ ca \end{array} \right] \left[ \begin{array}{c} ca \\ a \end{array} \right] \left[ \begin{array}{c} a \\ ab \end{array} \right] \left[ \begin{array}{c} abc \\ c \end{array} \right]$$

abcacabc

$$S_1 = \left\{ \left[ \frac{abc}{ab} \right], \left[ \frac{bc}{a} \right], \left[ \frac{cc}{b} \right], \left[ \frac{bb}{c} \right] \right\}$$

no solution b/c top string  
always longer

$$S_2 = \left\{ \left[ \frac{abc}{db} \right], \left[ \frac{bc}{af} \right], \left[ \frac{b}{df} \right] \right\}$$

no solution b/c no d, f  
is top row

Claim: PCP is undecidable  
(it is recognizable)

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Proof outline:

- prove MPCP undecidable
- show PCP  $\rightarrow$  MPCP

MPCP: like PCP but solution  
sequence must start with  
1st tile in list

To show MPCP undecidable,  
simulate a TM

Suppose  $M$  a TM & string  $w$   
Make set  $S$  of Tiles

a solution for  $S$  has top (bottom  
string

initial cuts for  $w$

#  $C_1$  #  $C_2$  # ..... #  $C_n$  ##

↑

configurations of  
 $M$

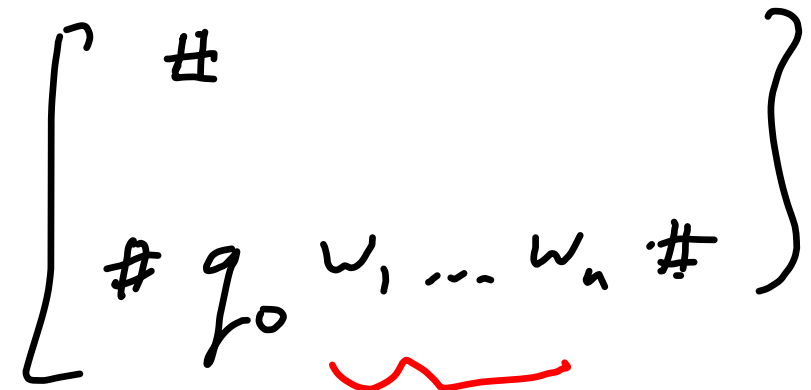
abc q defg

↑ almost

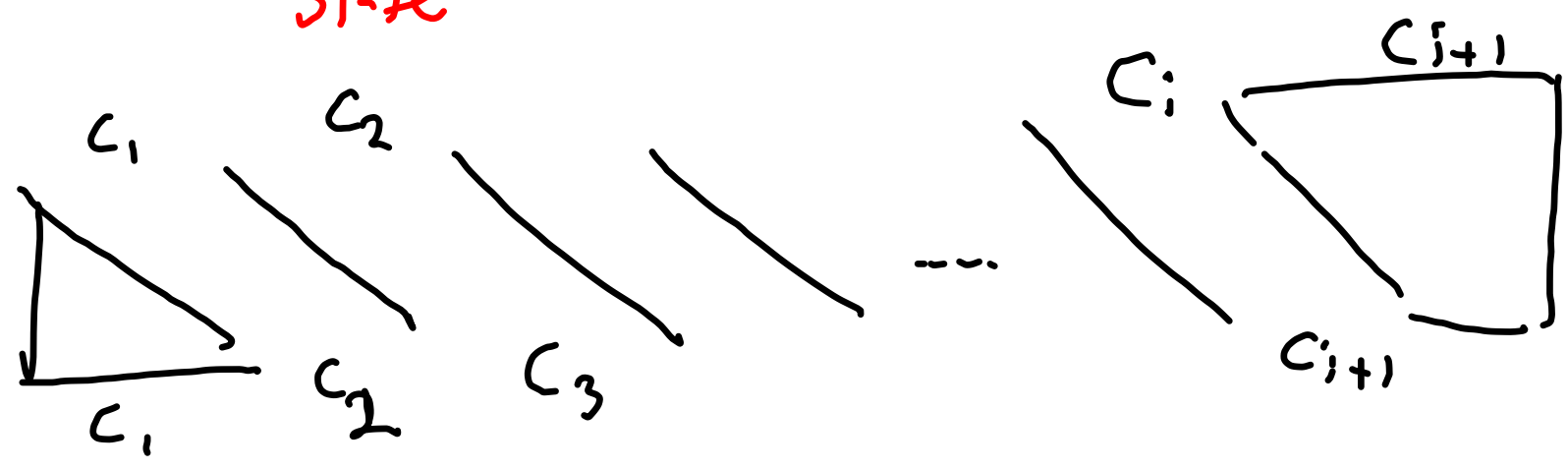
accepts  
configurations

Make set of tiles  $S$  s.t.  
 $\exists$  solution for  $S$  iff  
 $M$  accepts  $w$

start tile

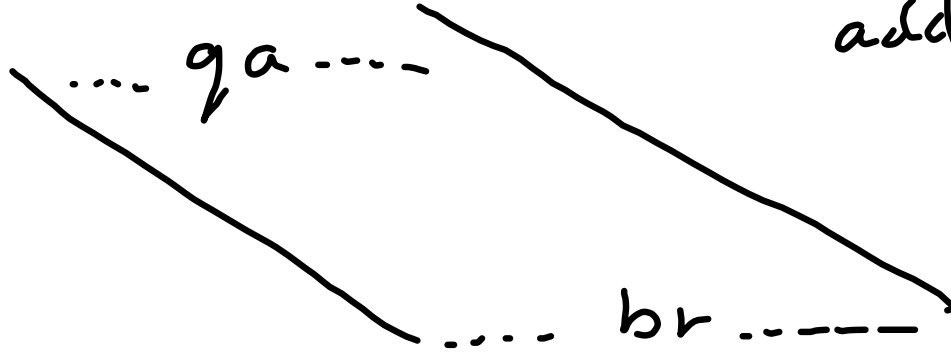


initial state      input string  $w$



actions

$c_i$



$c_j$

if  $\delta(q, a) = (r, b, R)$

add tile

$$\left[ \begin{array}{c} qa \\ \hline br \end{array} \right]$$

if  $\delta(q, a) = (r, b, L)$

add tiles

$$\left[ \begin{array}{c} cq a \\ \hline rcb \end{array} \right] \forall c \in \Sigma$$

# Copying files

add  $\begin{bmatrix} a \\ a \end{bmatrix}$   $\forall a \in \Sigma$   
to  $S$

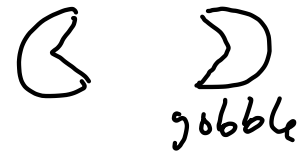
also add to  $S$   $\begin{bmatrix} \# \\ \# \end{bmatrix}$   $\begin{bmatrix} \# \\ \cup \# \end{bmatrix}$

end



add tile  $\begin{bmatrix} q_{accept} \# \\ \# \end{bmatrix}$

also add pacman transitions



$\begin{bmatrix} a & q_{accept} \\ q_{accept} & \end{bmatrix}$        $\begin{bmatrix} q_{accept} & a \\ & q_{accept} \end{bmatrix}$

$\forall a \in \Gamma$

If I have a set  $S$  of tiles  
for MPCP problem

Can I build equivalent set of  
tiles for PCP

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Suppose  $w = w_1 w_2 \dots w_n$

$*w = *w_1 *w_2 \dots *w_n$

$w* = w_1 *w_2 * \dots w_n *$

$*w* = *w_1 *w_2 * \dots *w_n *$

Suppose  
MPCP  $S = \left\{ \begin{bmatrix} t_1 \\ b_1 \end{bmatrix}, \begin{bmatrix} t_2 \\ b_2 \end{bmatrix}, \dots, \begin{bmatrix} t_n \\ b_n \end{bmatrix} \right\}$

must start sequence

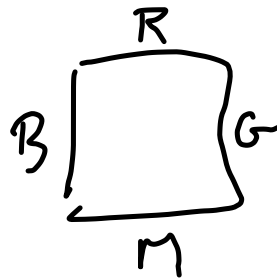
$$S' = \left\{ \begin{bmatrix} *t_1 \\ *b_1* \end{bmatrix}, \begin{bmatrix} *t_1 \\ b_1* \end{bmatrix}, \begin{bmatrix} *t_2 \\ b_2* \end{bmatrix}, \dots, \begin{bmatrix} *t_n \\ b_n* \end{bmatrix}, \begin{bmatrix} * \diamond \\ \diamond \end{bmatrix} \right\}$$

Wang tiles

1st aperiodic tiling

1966 Robert Berger

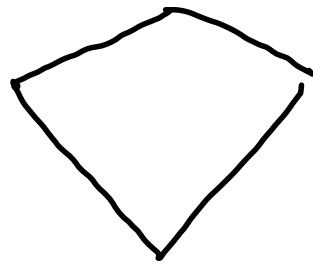
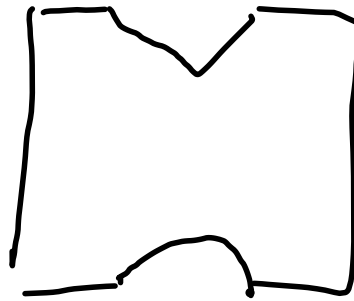
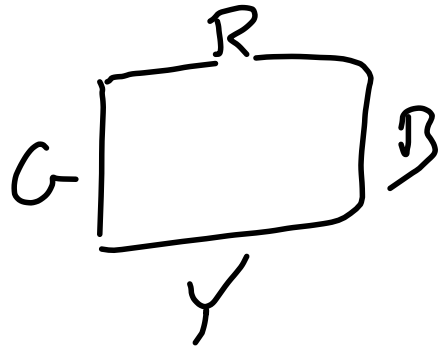
20,426 wang tiles



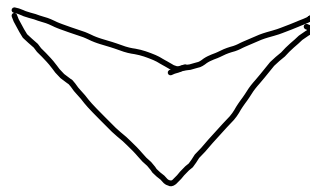
..... 1996

Karel Culik, II

13 wang tiles



kite



dart

Suppose That if a set  
of tiles  $S$  can  
tile whole plane,  
it can tile whole  
plane periodically.

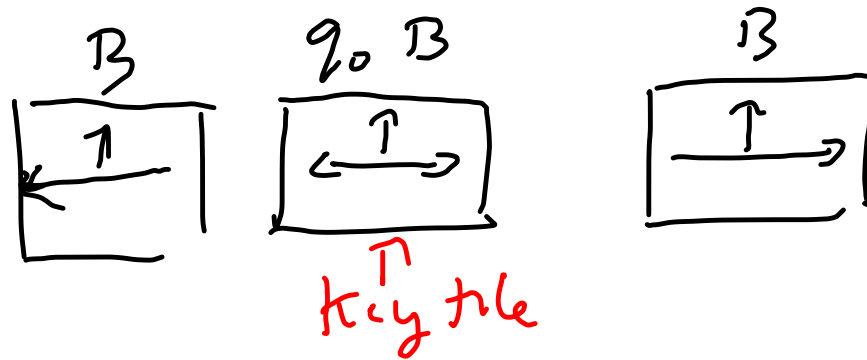
Conjecture by  
Hao Wang 1961

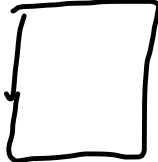
$\Rightarrow$  WRONG

# Tiles for "completion problem"

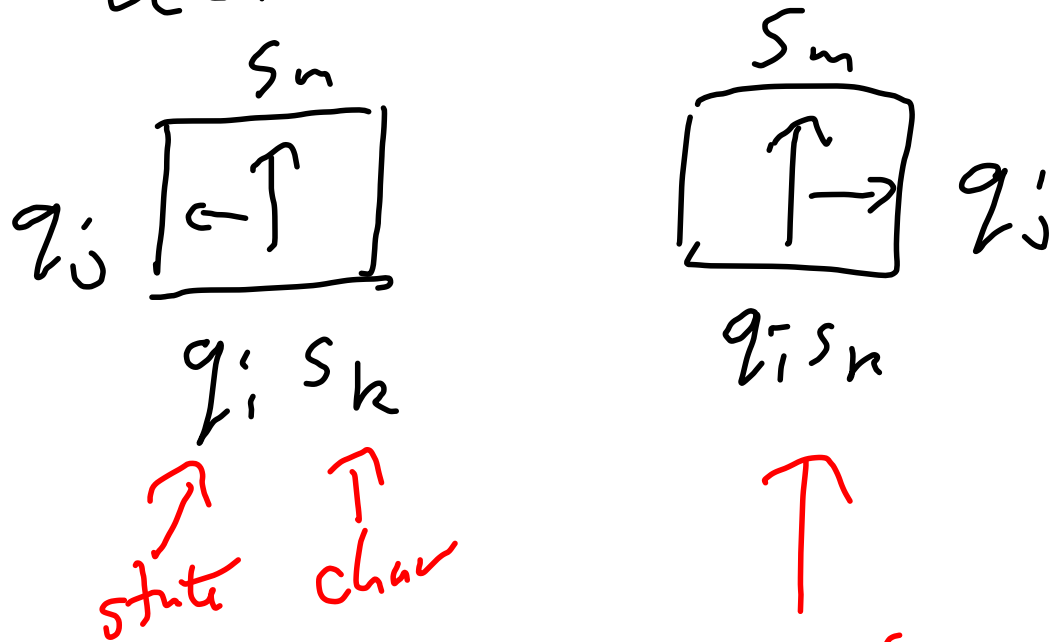
→ given part of tiling,  
can it be extended  
to  $\mathbb{R}^2$ ?

Start TM on blank input



also  bottom 1/2 plane

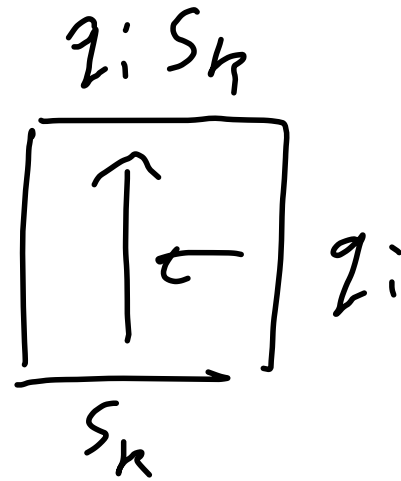
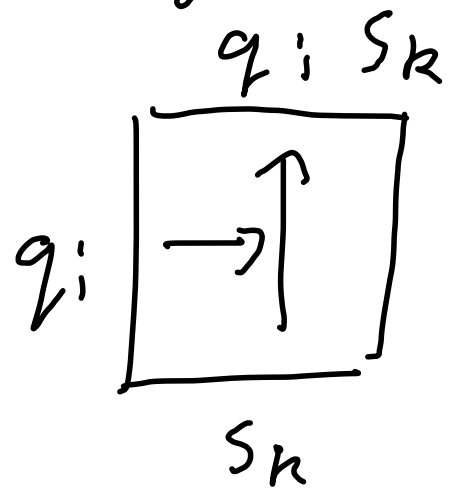
action files



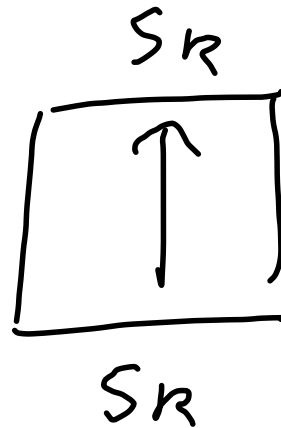
TM has

$$\delta(q_i, s_k) = (q_i, s_m, R)$$

merge tiles



copy tiles



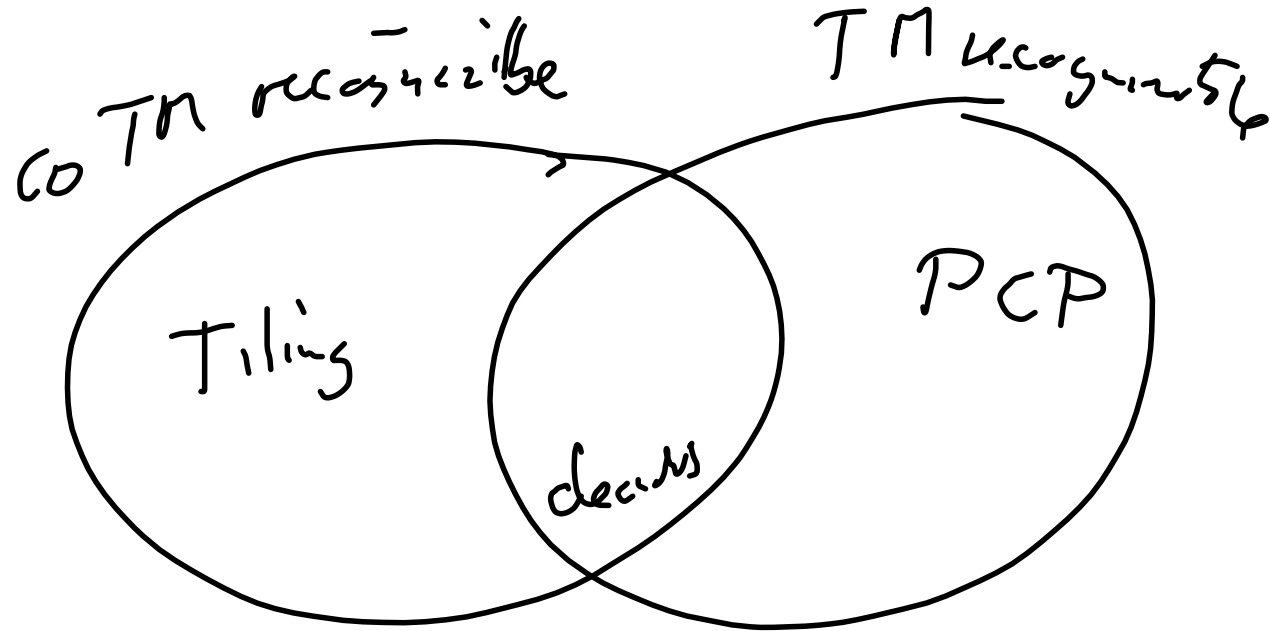
R. Robinson 1977

itsy-bitsy universal  
TM

(Minsky)

fixed set of 36 tiles

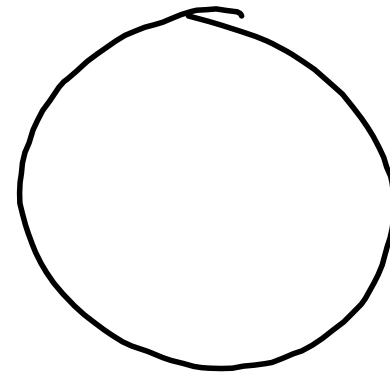
for which completion isn't  
decidable



# Knot diagrams

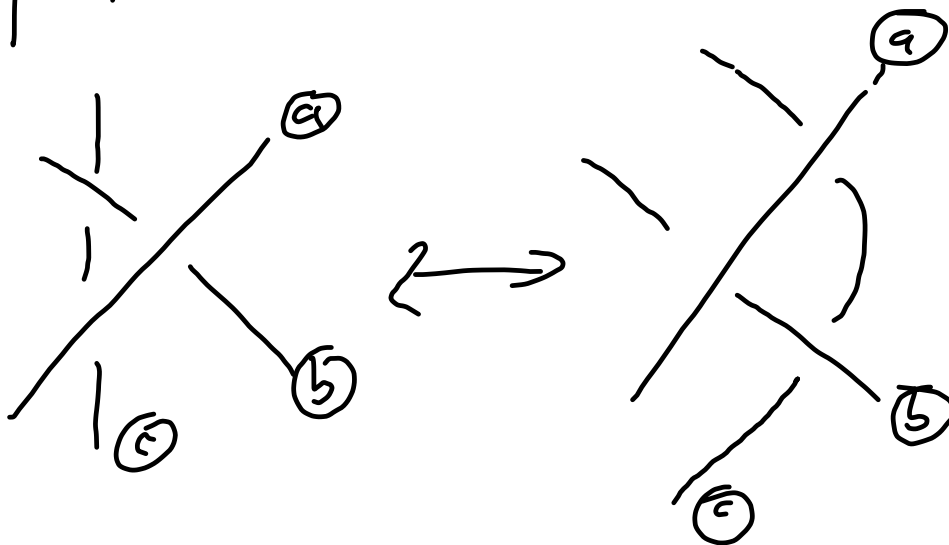
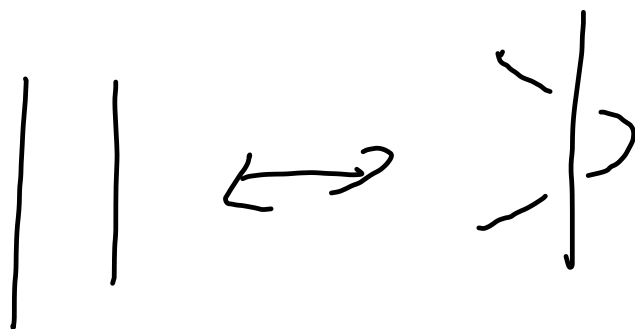
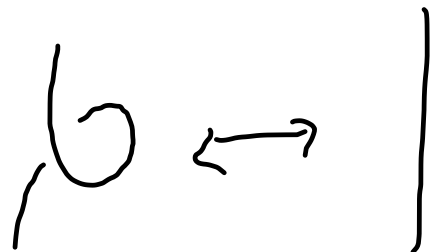


unknot



# Reidermeister moves

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1926/27 R. moves

1934 L. Goeritz

Sometimes must increase  
# of crossings

2001 Hass & Lagarias

$n$  crossings

$\leq 2^{cn}$  moves

$c \leq 10$  "