

a homomorphism h
 is a function that
 maps characters to
 strings

$$h: \Sigma \rightarrow \Gamma^*$$

$$\text{Ex: } \Sigma = \{a, b, c, \#\}$$

$$\Gamma = \{m, p, k\}$$

$$h: a \rightarrow m$$

$$b \rightarrow mpp$$

$$c \rightarrow \epsilon$$

$$\# \rightarrow k$$

$$h(c \# b a a) = \overset{\text{not there really}}{\epsilon} k \overset{b}{\overbrace{mpp}} mm$$

$$= k m p p m m$$

$$L = \{s^n t^n; n \geq 0\}$$

Claim

L can't be regular.

Suppose L is regular.

$$\text{Let } h: \begin{array}{l} s \rightarrow 0 \\ t \rightarrow 1 \end{array}$$

$$\text{Then } h(L) = \{0^n 1^n; n \geq 0\}$$

Since L is regular and
regular lgs are closed
under homomorphism

Then $h(L)$ must be regular

But we've shown that
 $h(L)$ isn't regular
#

$$L'' = \{s^n \# t^n : n \geq 0\}$$

$$h: s \rightarrow 0$$

$$t \rightarrow 1$$

$$\# \rightarrow \varepsilon \quad (\# \text{ disappears})$$

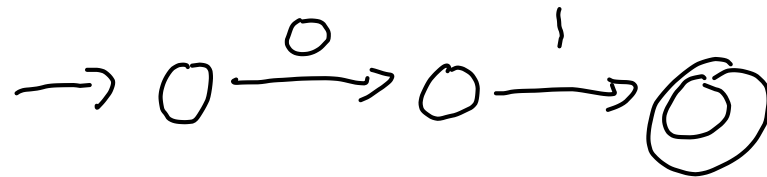
$$h(L'') = \{0^n 1^n : n \geq 0\}$$

input n NFA
need to build DFA

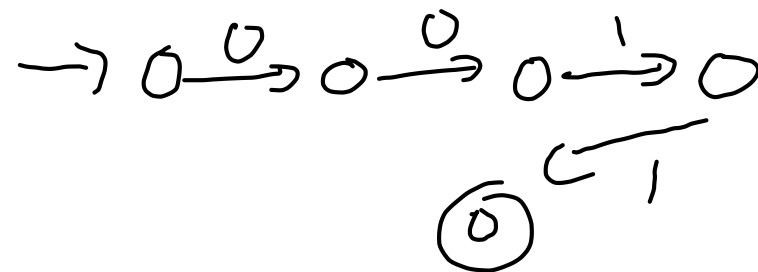
M_n recognises

$\{0^n, 1^n\}$ n fixed
constant

$n=1$



$n=2$



in general

$$M_n = (Q_n, \bar{\Sigma}, \delta, q_1, F_n)$$

$$Q_n = \{q_1, \dots, q_{2n+1}\}$$

$$F_n = \{q_{2n+1}\}$$

$$\delta(q_i, 0) = \{q_{i+1}\}$$

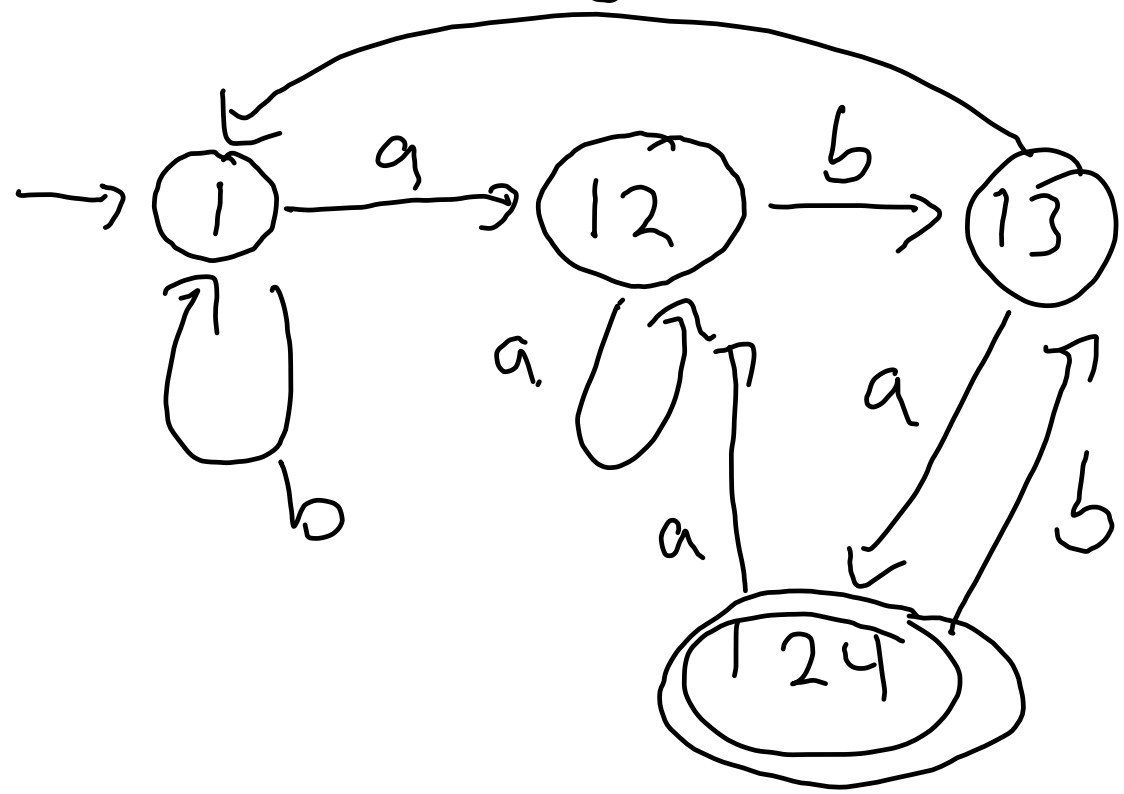
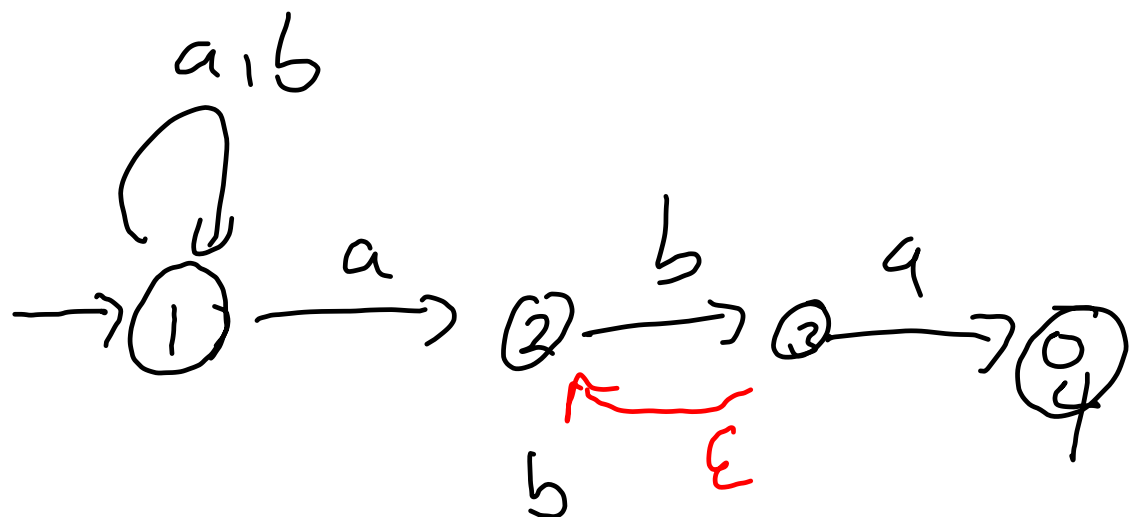
if $i \leq n$

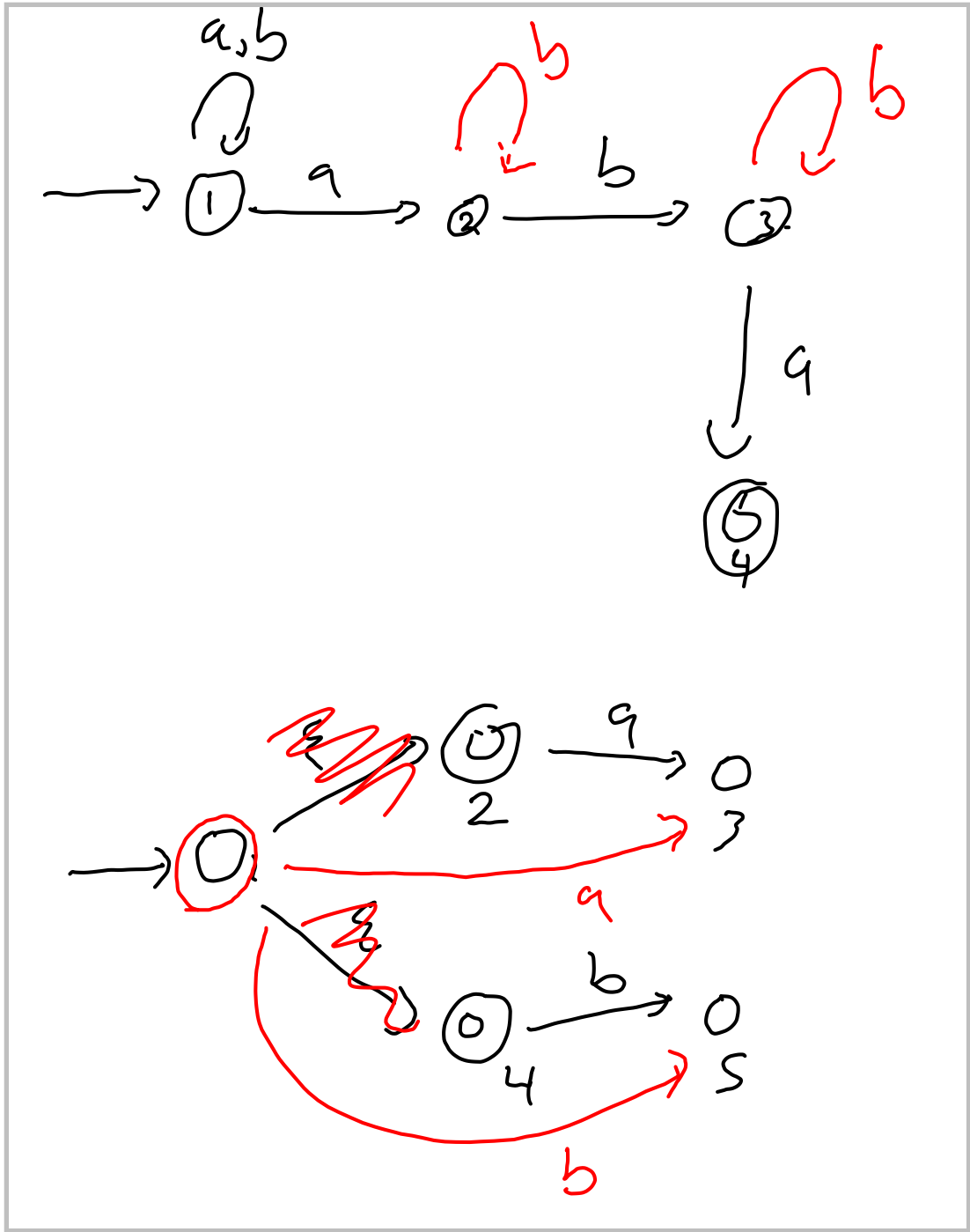
\emptyset otherwise

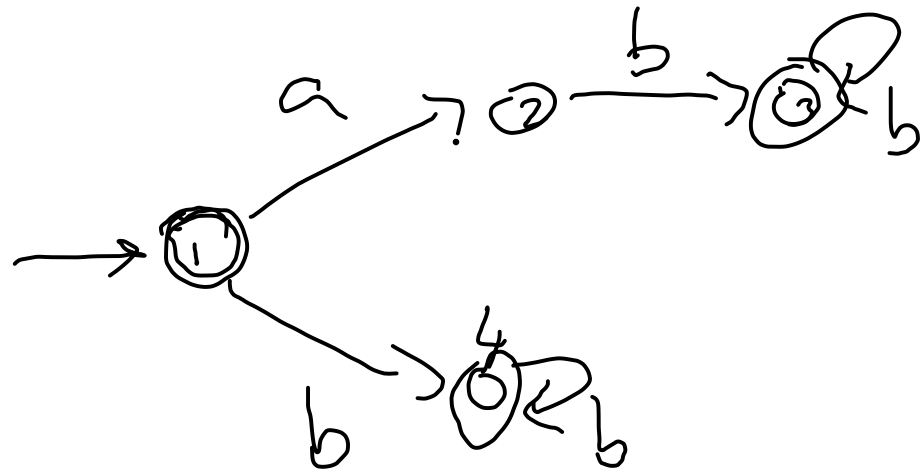
$$\delta(q_i, 1) = \{q_{i+1}\}$$

if $n < i \leq 2n$

\emptyset otherwise







1. $abb^* \cup bb^*$

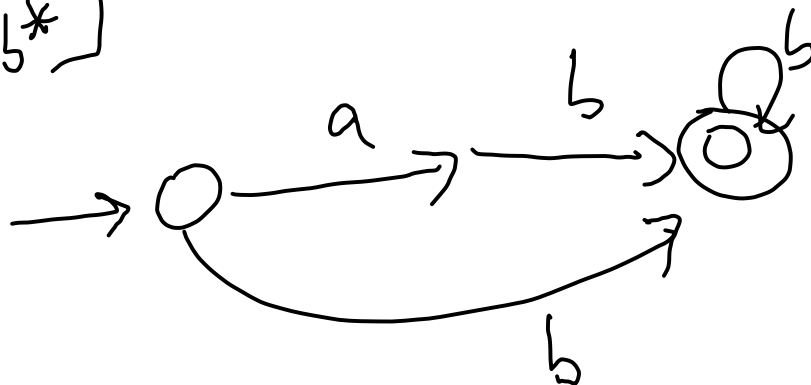
2. bb^* \rightarrow

3. b^* \rightarrow

4. b^* \rightarrow

abb^*

$\cup bb^*$



Generic construction

$$\text{NFA } M = (Q, \Sigma, q_0, \delta, F)$$



$$M' = (Q', \Sigma, q_0', \delta', \underbrace{F'}_{\{f_x\}})$$

NFA \rightarrow NFA with
only one
final state

$$Q' = Q \cup \{f_x\} \quad (\text{Assumes } f_x \notin Q)$$

$$\delta'(q, a) = \delta(q, a) \text{ if } \underline{q \in Q}$$

$$\delta'(q, \varepsilon) = f_x \text{ if } \underline{q \in F}$$

$$\delta'(f_x, a) = \emptyset$$